A fast-locking harmonic-free digital DLL for DDR3 and DDR4 SDRAMs

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Abstract: A new digital delay-locked loop (DLL) for DDR3/DDDR4 SDRAM is presented. The proposed digital DLL employs a new noise-tolerant triple (MSB-interval + binary + sequential) search algorithm for implementing a harmonic-free, fast-locking capability while retaining low jitter, low power performance, and a wide operating frequency range. The proposed DLL with duty-cycle correction is designed using a 38-nm CMOS process and occupies an active area of just 0.02 mm². The DLL operates over a frequency range of 0.3–2.0 GHz and achieves a peak-to-peak jitter of 7.78 ps and dissipates 3.48 mW from a 1.1 V supply at 1 GHz.

Keywords: delay-locked loop, DDR3, DDR4, SDRAM, harmonic-free, DLL

Classification: Integrated circuits

References

1 Introduction

Synchronous dynamic random access memory (SDRAM) has served as an important low-cost main memory solution for the personal computer (PC) and other cost-sensitive consumer electronics markets. Traditional SDRAM provides a maximum memory bus data rate of only 166 Mbps/pin with a maximum clock rate of 166 MHz. As processor speeds continue to increase, the PC becomes more reliant on low cost, higher bandwidth memory solutions. The double data rate (DDR) SDRAM was introduced to the public market to meet these requirements. DDR SDRAM, also called DDR1 SDRAM, continues to evolve and has been superseded by various DDR-x SDRAMs (i.e., DDR2, DDR3, and DDR4) that provide better performance and consume lower power. DDR3 was introduced in 2007 and DDR4 in 2013 [1]. Currently, the market is moving swiftly from DDR3 to DDR4 because of the lower power consumption and higher speed advantages of DDR4.

In order to achieve a memory bus data rate of over 400 Mbps/pin, DDR-x SDRAMs must incorporate an on-chip delay-locked loop (DLL) [1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12] that can eliminate skew problems and achieve higher timing margin at high frequencies. To design a DLL that can support both DDR3 and DDR4 specifications at the same time [13, 14], the DLL should be locked within 512 clock cycles and operate over a frequency range from 300 MHz to 1.6 GHz using an internal supply voltage of less than 1.2 V. Also, the DLL must be capable of correcting the duty cycle of the distorted input clock so that the data-valid window (tDV) could be widened [2].

Currently, most DDR3/DDR4 SDRAMs use a digital DLL [1, 2, 3, 4, 10, 11]. One of the reasons for using digital architectures is because DDR3/DDR4 SDRAMs require fast recovery times for various power mode transitions. To
achieve a fast locking time, a successive approximation register (SAR)-based binary search algorithm was adopted in DLL designs [5, 6, 7, 8, 9]. However, this introduced harmonic lock problem [6, 7, 12]. Harmonic locking may occur when the most significant bit (MSB) of the SAR code is changed in the beginning of the SAR operation, since the delay of the DLL is increased to 50% of the total delay line. To eliminate the harmonic locking problem in a SAR-based digital DLL, a variable SAR (VSAR) algorithm was introduced [7]. However, [7] requires a very complex and timing sensitive fail-to-lock detection circuit which could easily cause logic failures in the presence of noise in the power supply or ground.

In this paper, a digital DLL for DDR3/DDR4 SDRAMs is proposed which relies on a new noise-tolerant triple (MSB-interval + binary + sequential) search algorithm for achieving a harmonic-free, fast-locking capability while maintaining low jitter, low power consumption, and a wide operating frequency range.

### 2 Circuit design

Fig. 1(a) shows a block diagram of the proposed digital DLL for DDR3/DDR4 SDRAMs. It consists of a clock buffer, a digitally-controlled delay line (DCDL), a duty-cycle corrector (DCC), a DCDL controller, a divider, a phase detector (PD), a clock tree, and a replica delay line (RDL). The DCDL comprises a coarse delay line

![Block Diagram of Proposed Digital DLL](image)

*Fig. 1. (a) Proposed digital DLL architecture (b) Flow chart of the proposed triple search algorithm with three operating modes*
The DCDL controller consists of a control logic, a 10-bit multi-mode register (MMR) and two 5-to-32 thermometer decoders. The clock tree, which introduces a delay of t2, is a clock distribution network connected to the output drivers (DQs) of the SDRAM. The RDL is a dummy delay that replicates, with a smaller area and lower power overhead, the propagation delay of the clock tree. To attain phase aligned locking between the CLK/CLKB and the CLKDQ without clock skews, the delay of the RDL (t3) is given by the equation \( t3 = t1 + t2 \), where t1 is the CLK buffer delay. Many of the previously reported fast-locking DLL schemes cannot be implemented in DDR-x SDRAMs due to the existence of the RDL which affects the locking time and the harmonic locking issue.

The PD takes CLKIN and CLKFB as inputs and generates a Comp signal that is used by the DCDL controller to adjust the DCDL delay. M[4:0], the five most-significant-bits (MSBs) of the 10-bit MMR, are converted to the thermometer codes, C[31:0]/Cc[31:0], by the 5-to-32 decoder. The codes C[31:0]/Cc[31:0] are used for controlling the CDL of the DCDL. L[4:0], the five least-significant bits (LSBs) of the MMR, are converted to the thermometer codes, F[31:0]/Fb[31:0]. The codes F[31:0]/Fb[31:0] are used for adjusting the magnitude of the delay in the FDL.

This digital DLL architecture has three operating modes (MSB, SAR, and Counter mode) and utilizes a triple (MSB-interval + binary + sequential) search algorithm for achieving fast locking without harmonic locking, as shown in Fig. 1(b).

The MSB mode relies on eight evenly spaced delay intervals (1’st~8th interval) to control the 5-bit CDL, which provides high noise tolerance in the DCDL controller design.

Fig. 2 shows the proposed CDL which is made up of 32 cascaded digital delay elements (DEs). The DE consists of four NAND gates [7] and the delay step in a
DE is equivalent to two NAND gates. Therefore the maximum total delay of the CDL is \(32 \times t_D\), where \(t_D\) is the propagation delay of the DE.

The delay of the CDL is controlled by the 32-bit thermometer code \(C[31:0]/C_b[31:0]\). As the operating frequency decreases, the required number of activated DEs increases. The fine delay line (FDL) is based on feedback delay element (FDE)-based inverters [9] and the tunable delay range of the FDL is equal to one DE delay, \(t_D\). Since the FDL delay is incremented or decremented by the 32-bit \(F[31:0]/F_b[31:0]\) code, the delay resolution of the proposed digital DLL is about \(t_D/2^5 = (135 \text{ ps})/32 = 4.2 \text{ ps}\).

The propagation delay of the DLL can be represented as follows

\[
t_{\text{propagation delay}} = t_{\text{variable}} + t_{\text{fixed}} = N \times t_{\text{CK}}
\]

where \(t_{\text{variable}}\) is the tunable delay of the DCDL, \(t_{\text{fixed}}\) is the initial fixed delay of the DLL when the \(t_{\text{variable}}\) equals to zero, and \(t_{\text{CK}}\) is the cycle time of the input clock. For ideal phase lock locking without the harmonic lock problem, \(N\) needs to be one. Therefore, \(t_{\text{variable}}\) can be represented as follows.

\[
t_{\text{variable}} = 1 \times t_{\text{CK}} - t_{\text{fixed}}
\]

Since the CDL consists of 32 DEs and the tunable delay range of the FDL is equal to one \(t_D\), \(t_{\text{variable}}\) becomes as follows.

\[
t_{\text{variable}} = (32 + 1) \times t_D = 33 \times t_D
\]

In order to achieve harmonic-free operation at the minimum frequency of 0.3 GHz (\(t_{\text{CK}} = 3.33 \text{ ns}\)) with any values of \(t_{\text{fixed}}\) and \(t_3\), the maximum variable delay of the DCDL (\(t_{\text{variable, max}}\)) should be at least 3.33 ns. Therefore, the proper value of \(t_D\) should be larger than 100 ps (= \(t_{\text{variable, max}}/33 = 3.33 \text{ ns}/33\)) even in the fastest process corners. Consequently, by choosing a typical \(t_D\) of 135 ps, the proposed DLL can achieve harmonic-free wide-range operation from 300 MHz to 2.0 GHz.

Referring to Fig. 1(b) and Fig. 2, the DLL starts in the MSB mode with \(M[4:0] = [00010]\), which enables the two DEs (#1 and #2) of the CDL and subsequently the MSB-interval search algorithm is executed. Therefore, at the beginning, the DLL starts in the 1’st interval period and then moves to the 2’nd interval period regardless of the Comp signal by enabling four DEs (#1–#4) of the CDL. Starting from the 2’nd interval period, the DLL moves to the next interval period if the Comp signal is logic high. For instance, \(M[4:0] = [11100]\) represents the 8’th interval period, in which twenty eight DEs (#1–#28) are enabled in the CDL. When the Comp signal is changed from logic high to low, the MSB mode is completed and the DLL enters the SAR mode and the binary search algorithm begins.

Fig. 3(a) illustrates the triple search locking process along with the DCDL control bits and the three operating modes. The MSB mode is used to set the DLL output clock near the locking point within eight CLKCTRL cycles, where CLKCTRL is an output of the divide-by-N divider with \(N = 4\). Thus the delay range of the DCDL is separated by eight interval periods, which produces a maximum delay step change of only \(4 \times t_D\) in the MSB mode. When the MSB mode is completed, the DLL enters the SAR mode and the binary search algorithm is applied to the MMR.
Fig. 3(b) shows the 10-bit MMR. In order to increase the DCDL delay rapidly without incurring any harmonic locking problem, only the 3 MSB bits, M[4:2], are controlled in the MSB mode. Then the rest 7 LSB bits are controlled in the SAR mode for binary search. The SAR-based binary search requires a maximum 7 $\text{CLK}_{\text{REG}}$ cycles, where $\text{CLK}_{\text{REG}}$ is an output of the divide-by-M divider with $M = 8$.

After the binary search is completed in the SAR mode, the 10-bit MMR is transformed into a 10-bit counter and the DLL starts the sequential search in the Counter mode, maintaining a closed loop to track process, voltage, and temperature (PVT) variations. As a consequence, the use of triple search algorithm results in a relatively fast locking time with no harmonic locking problems. The worst case locking time of the proposed DLL is $88 \times (8 \times \text{CLK}_{\text{CTRL}} + 7 \times \text{CLK}_{\text{REG}} = 8 \times 4 + 7 \times 8)$ input clock cycles.

Fig. 4(a) and (b) show the proposed control logic and the 10-bit MMR of the DCDL controller. Fig. 4(c) shows the truth table in the MSB mode and describes the relationship between $K[3:0]$ and $M[4:1]$ and the number of active DEs. In the MSB mode, $M[4:2]$, the three MSBs of the 10-bit MMR and representing the eight interval sequences from 000 to 111, is the same as $K[3:1]$. Also, the bits $M[1:0]$ are set to [00] and $L[4:0]$, the five LSBs of the MMR, are also set to [00000]. Therefore, $M[4:0]$, the five MSBs, are converted to the thermometer codes, $C[31:0]/C_b[31:0]$, using the 5-to-32 decoder. The signal $C[31:0]/C_b[31:0]$ is used for controlling the number of active delay elements of the CDL. To avoid harmonic locking problems, the number of active DEs in the MSB mode is increased almost linearly (from 2 to 28) with a maximum delay step change of only $4 \times t_D$. 
### 3 Experiment results

The proposed digital DLL was designed in a 38-nm Powerchip CMOS process. Fig. 5(a) shows a chip layout which has an active area of 153 µm × 137 µm (≈ 0.02 mm²). To minimize process variations and device mismatch problems, 70 nm was used as the minimum channel length for the simulation and layout in this design. Fig. 5(b) shows a layout of the 2 Gb DDR3 SDRAM which incorporates the proposed DLL.

The timing diagram of Fig. 6 shows the simulated locking process of the proposed digital DLL at 300 MHz. The DLL starts with \( M[4:0] = [00010] \) and \( L[4:0] = [00000] \) (≈ 64 in decimal). It can be seen that the MSB mode is completed at 481 ns with \( M[4:0] = [11000] \) and \( L[4:0] = [00000] \) (≈ 768 in decimal). In SAR mode, \( M[4:0] = [10110] \) and \( L[4:0] = [10001] \) (≈ 721 in decimal). In Counter mode, the input clock (\( CLK_{IN} \)) and output clock (\( CLK_{OUT} \)) of the DLL are precisely aligned with each other with no clock skews. It is assumed that the clock buffer and the replica path are ideal and therefore they are omitted in this simulation.

Fig. 7 shows the simulated peak-to-peak (p-p) jitter of the output clock. The proposed DLL achieves a simulated p-p jitter of 12.03 ps and 9.71 ps at 300 MHz.

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**Table 1:**

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<thead>
<tr>
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<tbody>
<tr>
<td>1st internal</td>
<td>0001</td>
<td>0001</td>
<td>2</td>
</tr>
<tr>
<td>2nd internal</td>
<td>0010</td>
<td>0010</td>
<td>4</td>
</tr>
<tr>
<td>3rd internal</td>
<td>0100</td>
<td>0100</td>
<td>8</td>
</tr>
<tr>
<td>F1 interval</td>
<td>1100</td>
<td>1100</td>
<td>12</td>
</tr>
<tr>
<td>F2 interval</td>
<td>1000</td>
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<tr>
<td>F4 interval</td>
<td>1110</td>
<td>1110</td>
<td>24</td>
</tr>
<tr>
<td>F5 interval</td>
<td>1110</td>
<td>1110</td>
<td>28</td>
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</tbody>
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**Fig. 4.** Proposed DCDL controller (a) control logic (b) 10-bit multi-mode register (MMR) (c) truth table in the MSB mode
and 2.0 GHz, respectively. Fig. 8 shows the simulated p-p jitter of the output clock when there exists a 1 GHz supply noise of 50 mV. Although the p-p jitter has been increased to 46 ps with this injected supply noise, the DLL shows good noise-

Fig. 5. Chip layout of the (a) proposed digital DLL (b) DDR3 SDRAM with the proposed DLL.

Fig. 6. Simulated locking process of the proposed digital DLL.

Fig. 7. Simulated peak-to-peak jitter (a) 300 MHz (b) 2 GHz.
tolerant operation. The DLL operates over a wide frequency range of 0.3–2.0 GHz, which encompasses the operating frequency ranges of both DDR3 and DDR4. The proposed DCC-equipped DLL dissipates 3.48 mW at 1 GHz from a supply voltage of 1.1 V.

In DDR memory systems, the input clock of the DLL may have a duty-cycle distortion (DCD). Therefore, the DLL should be able to lock properly regardless of the input clock’s duty-cycle variation. In this design, the PD compares only the rising edges of the input and feedback clocks, and the DCDL controller has no relationship with the input clock’s duty-cycle ratio. Therefore, the phase locking operation of the proposed DLL is immune to DCD. Since the output of the DLL clock signal can be distorted due to device mismatches, the proposed DLL adopts a digital DCC [9] to improve performance at high frequencies. The digital DCC achieves a fast locking time of less than 24 input clock cycles and it can be turned off during the power-down mode. The DCC achieves a duty-cycle correction range of 30–70% at 1.0 GHz. Fig. 9 shows the simulated input and output clocks of the proposed DCC-equipped DLL with distorted input clock duty-cycles. As shown in Fig. 9(a), the DLL achieves a corrected output clock duty-cycle of 50.38% from a 55% input duty-cycle at 300 MHz. Fig. 9(b) shows an output duty-cycle of 50.48% from a 60% input duty-cycle at 2 GHz.

A performance comparison between the proposed digital DLL and other state-of-the-art DDR3/DDR4 digital DLLs is given in Table I.
This Letter presents a new digital DLL architecture which is capable of duty-cycle correction and harmonic-free, fast-locking in DDR3 and DDR4 SDRAMs. The proposed DLL achieves these capabilities by adopting a new noise-tolerant triple (MSB-interval + binary + sequential) search algorithm. Implemented in a 38-nm Powerchip CMOS DRAM process, the proposed DLL operates over a frequency range of 0.3–2.0 GHz, achieves a peak-to-peak jitter of 7.78 ps, and dissipates 3.48 mW from a 1.1 V supply at 1 GHz. And the proposed DCC-equipped DLL occupies an active area of just 0.02 mm².

Table I. Performance Summary and Comparison with state-of-the-art DDR3/DDR4 digital DLLs

<table>
<thead>
<tr>
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<tr>
<td>Process</td>
<td>65 nm</td>
<td>45 nm</td>
<td>0.25 µm</td>
<td>0.13 µm</td>
<td>0.18 µm</td>
<td>38 nm</td>
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<td>Architecture</td>
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<td>Digital</td>
<td>Digital</td>
<td>Digital</td>
<td>Digital</td>
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</tr>
<tr>
<td>Supply (V)</td>
<td>1.0–1.5</td>
<td>1.1</td>
<td>1.1</td>
<td>1.2</td>
<td>1.8</td>
<td>1.1</td>
</tr>
<tr>
<td>Frequency range (GHz)</td>
<td>0.12–2</td>
<td>0.4–0.8</td>
<td>0.1</td>
<td>0.03–1</td>
<td>0.04–0.55</td>
<td>0.3–2.0</td>
</tr>
<tr>
<td>Pk-to-pk jitter (ps)</td>
<td>14 @2 GHz</td>
<td>20.4 @800 MHz</td>
<td>95 @100 MHz</td>
<td>30 @1 GHz</td>
<td>16.9 @200 MHz</td>
<td>9.71 @2 Hz</td>
</tr>
<tr>
<td>Duty-cycle Correction</td>
<td>O</td>
<td>O</td>
<td>X</td>
<td>X</td>
<td>O</td>
<td>O</td>
</tr>
<tr>
<td>Locking time</td>
<td>-</td>
<td>-</td>
<td>30 cycles</td>
<td>42 cycles</td>
<td>134 cycles</td>
<td>88 cycles</td>
</tr>
<tr>
<td>Power</td>
<td>6.6 mW @1.2 V, 2 GHz</td>
<td>3.3 mW @800 MHz</td>
<td>3.3 mW @100 MHz</td>
<td>3.6 mW @1 GHz</td>
<td>9 mW @200 MHz</td>
<td>3.48 mW @1 GHz</td>
</tr>
<tr>
<td>Active area (mm²)</td>
<td>0.059</td>
<td>0.01</td>
<td>0.136</td>
<td>0.2</td>
<td>0.2</td>
<td>0.02</td>
</tr>
</tbody>
</table>

4 Conclusion

This Letter presents a new digital DLL architecture which is capable of duty-cycle correction and harmonic-free, fast-locking in DDR3 and DDR4 SDRAMs. The proposed DLL achieves these capabilities by adopting a new noise-tolerant triple (MSB-interval + binary + sequential) search algorithm. Implemented in a 38-nm Powerchip CMOS DRAM process, the proposed DLL operates over a frequency range of 0.3–2.0 GHz, achieves a peak-to-peak jitter of 7.78 ps, and dissipates 3.48 mW from a 1.1 V supply at 1 GHz. And the proposed DCC-equipped DLL occupies an active area of just 0.02 mm².

Acknowledgments

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