Distributed Estimation over Delayed Sensor Network with Scalable Communication

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SUMMARY This paper proposes a distributed delay-compensated observer for a wireless sensor network with delay. Each node of the sensor network aggregates data from the other nodes and sends the aggregated data to the neighbor nodes. In this communication, each node also compensates communication delays among the neighbor nodes. Therefore, all of the nodes can synchronize their sensor measurements using scalable and local communication in real-time. All of the nodes estimate the state variables of a system simultaneously. The observer in each node is similar to the delay-compensated observer with multi-sensor delays proposed by Watanabe et al. Convergence rates for the proposed observer can be arbitrarily designed regardless of the communication delays. The effectiveness of the proposed method is verified by a numerical simulation.

key words: sensor network, data aggregation, communication delay, distributed estimation

1. Introduction

Sensor networks with numerous sensors have attracted much attention because the development of micro-electro-mechanical systems (MEMS) has improved the performance of compact sensors and communication elements [1]. Many sensors can realize a wide-range of complex observations in large-scale systems. Redundant sensors improve the accuracy and robustness of an observation, and enable a fault-tolerant observation. A flexible sensing system can be realized by connecting sensors via a wireless network. Such a sensor network can be used in various applications, including area surveillance or the active monitoring of forests and agricultural lands.

There have been many studies on the applications of sensor networks [2]–[9]. Distributed estimation methods have also been proposed to reduce wasteful communication paths. Olfati-Saber et al. [10]–[13] proposed a distributed Kalman filter based on a consensus filter. The consensus filter is an application of consensus controls, and provides the average consensus of all the sensors included in a network. The consensus filter can calculate the consensus value through communications between adjacent nodes. An estimate of the distributed Kalman filter is obtained from the original Kalman filter using the consensus value [14]. Olfati-Saber et al. also proposed a Kalman-consensus filter, which executes the Kalman filtering and consensus calculation simultaneously. A gossip algorithm is also a distributed consensus algorithm for sensor networks [15]. In the gossip algorithm, each node selects the data sent from the other nodes at random. This reduces the communication traffic of the sensor network, and allows this algorithm to obtain a consensus value.

These studies focus on the communication efficiency. The communication delay is another problem of a sensor network with delay. In particular, the communication delays in a large wireless network cannot be ignored. However, the distributed estimation methods discussed in previous studies do not consider the communication delay. Delay compensated observers, which do not assume a network structure, have been proposed in [16]–[18]. In particular, Watanabe et al. [16] and Tsubakino et al. [18] considered the case where the output vector includes multiple delays. The delays included in the network are non-uniform, because they depend on the communication paths. The design of Watanabe’s observer resolves itself into a finite pole assignment problem.

In this paper, we bring a network structure to Watanabe’s delay-compensated observer. A distributed estimation method with delay compensation is proposed here. Watanabe defined an output vector, where each element of the vector is a sum of multiple measurements of a physical quantity with different delays. The output form of Watanabe’s observer is useful in aggregating the observed values using distributed data aggregation methods. Data aggregation methods for a sensor network are proposed in [19]. In this paper, we introduce the delay compensation proposed in [16] for tree-based data aggregation. All of the observed values of the sensor network are aggregated through communications between the neighbor-node pairs in a tree network. The observed values of the sensor network are aggregated by the tree-based communication at the root node. The communication delays are compensated by the memory of the input stored by each node. An observer of the root node can estimate the state from the aggregated data at its own node. We also propose an intercommunication protocol to aggregate the observed values in all the nodes, which is based on the fact that each node of a tree network can be a root node. Finally, the distributed observer can estimate the states at all the nodes. A dimension of the communication data among the nodes corresponds to a dimension of the output vector. Because the dimension of the output vector is independent of the number of nodes, the proposed communication law in this paper is scalable.
2. Problem Formulation

Figure 1 shows a system with distributed controllers. The dynamics of the plant in Fig. 1 can be expressed by

$$\dot{x} = Ax + Bu,$$

where $x \in \mathbb{R}^n$ is the state system, and $u \in \mathbb{R}^m$ represents the input values of the distributed controllers. Some states of system (1) can be measured by sensors as

$$y = Cx,$$

where $y \in \mathbb{R}^p$ is an output vector that consists of all of the redundant raw measurements, and therefore rank $C$ may be less than $p$, and $p$ may be larger than $n$. For the scalability of the network communication with respect to the number of sensors, this paper considers aggregations of sensor measurements. To reduce a communication amount, we utilize a data aggregation of the measurements. The raw output $y$ is aggregated into a $q$-dimensional vector $y_{agr}$ as

$$y_{agr} = Fy = C_{agr}x,$$

where $C_{agr} = FC$. Note that $q < p$ because the data aggregation reduce the dimension of original output $y$. The designs of $F$ and $q$ are important because these affect the performances of the state estimation or control, as indicated in [20]. However, this paper focuses on other problems, and we assume that $F$ is given.

There are two networks in Fig. 1, i.e., a sensor network and an actuator network that share the same node set with $N$ elements. Each node may have the sensors and a controller. Each controller collects the observed information from the other nodes via the sensor network to estimate the state $x$, and calculates a part of the input elements from the estimated value. By renumbering the elements of $y$, we can decompose matrix $F$ as

$$F = \begin{bmatrix} F_1 & \cdots & F_N \end{bmatrix},$$

where $F_i$ corresponds to the output of the $i$-th node. Therefore, the output of node $i$, which is mapped to the aggregated output space, is defined as

$$y_1 = \begin{bmatrix} F_1 & 0 & \cdots & 0 \end{bmatrix}y = C_1x,$$

$$\vdots$$

$$y_N = \begin{bmatrix} 0 & \cdots & 0 & F_N \end{bmatrix}y = C_Nx,$$

where $C_i = \begin{bmatrix} 0 & \cdots & 0 & F_i & 0 & \cdots & 0 \end{bmatrix}$. If the current outputs $y_i$ can be obtained with no transmission delay, the aggregated output coincides with $y_{agr}$, i.e.

$$y_{agr} = \sum_{i=1}^{N} y_i.$$

Each node needs all of the input values to estimate the state. In this paper, it is assumed that the dimension of the input is smaller than that of the output. Each node sends the input values via a high-speed actuator network with a limited capacity. On the other hand, the observed values of each of the nodes are sent via a sensor network with sufficient bandwidth but low communication speed. Therefore, the communication delay in the transmission of the input values is sufficiently smaller than that for the observed values. We have ignored the communication delay in the broadcast of the input values, and it is assumed that all the nodes can obtain the input instantly.

We represent the sensor network by an undirected graph. The set of nodes is denoted by $V := \{1, 2, \ldots, N\}$, and the set of edges is denoted by $E \subseteq V \times V$. By using $V$ and $E$, the undirected graph is expressed by $G(V, E)$. In this paper, it is assumed that the graph $G$ is a connected graph. For the connected graph $G$, there always exists at least one tree that is a subgraph of $G$ and includes all the nodes of $V$. This tree is denoted by $T(V, E)$, where $E$ satisfies $\hat{E} \subseteq E$. Node $i$ can mutually communicate with the neighbor nodes. The set of neighbors of node $i$ connected by $T$ is denoted by $J_{i} := \{j; (i, j) \in \hat{E}\}$. Once a root node of the undirected tree is chosen, the parent node and child nodes of each node are automatically determined. A set of the child nodes of node $i$ is denoted by $H_{ir}$, where the $r$-th node is chosen as the root node. From the definition $H_{ir}$ and $J_{i}$, $H_{ir} \subseteq J_{i}$. If $r = i$, $H_{ir} = J_{i}$. Otherwise, $p_{ir} = J_{i} \setminus H_{ir}$, where $p_{ir}$ is a parent node of node $i$ when $r$ is the root. For example, the set of child nodes of node 4 in Fig. 2 is

$$H_{2r} = \begin{cases} \{1, 4, 5\} & \text{if } r = 2 \\ \{1, 5\} & \text{if } r = 4, 8, 9 \\ \{1, 4\} & \text{if } r = 5 \\ \{4, 5\} & \text{if } r = 1, 3, 6, 7 \end{cases}.$$
\[ D_{42} + D_{21}. \]

**Remark 1.** In many cases, it is assumed that \( D_{ij} = D_{ji} \), which is a natural assumption. However, the method proposed in this paper does not need this assumption.

**Remark 2.** The sensor network may have relay nodes. A relay node does not have a sensor but can communicate with the other nodes. The output matrix of the relay node is \( C_i = 0 \). Several systems can be realized using relay nodes. For example, a system with a single centralized controller can be expressed by a network that includes a controller and has a relay node as the root node.

Main problems in this paper is design a consensus communication law for distributed observer over delayed sensor networks. For the communication and estimation, node \( i \) has \( \Lambda_i := \{A, B, \{D_{ij}; j \in J_i\}, C_i, \tilde{C}_i, \{\tilde{C}_{ij}; j \in J_i\}\} \) in own memory, where \( \tilde{C}_i \) and \( \tilde{C}_{ij} \) will be defined in Sects. 3.2 and 3.3. From the actuator network, the all nodes can obtain \( u(t) \) in real-time and store the history of input \( u(\tau) \), \( \tau \in [t, t - D_{i,\max}] \), where \( D_{i,\max} \) is the maximum delay expressed by \( D_{i,\max} = \max_{j \neq j} D_{ij} \). To simplify the problem, we divide the problem into the following two parts. In Problem 1, the root \( r \) is fixed and each node sends a message to own parent node. The message sent to the parent node by node \( i \) is \( (\hat{y}_i(t), \hat{\Sigma}_i(t)) \), where \( \hat{y}_i(t) \) is an aggregate measurement and \( \hat{\Sigma}_i(t) \) is a compensation value for communication delays. Then, only the root estimates the state as follows:

**Problem 1.** Assume that node \( r \) is fixed as the root node of \( T \) and \( (A, \tilde{C}_r) \) is the observable pair. Given \( \Lambda_i \) for all \( i \in V \). Then, find the communication law

\[
\hat{y}_i(t) = \hat{y}_i(y_i, \{\hat{y}_j(t - D_{ij}); j \in J_i\}), \\
\hat{\Sigma}_i(t) = \hat{h}_i(\hat{\Sigma}_i(t - D_{ij}); j \in J_i), \\
\{u(\tau); t - D_{i,\max} \leq \tau \leq t\},
\]

and the observer in the root

\[
\hat{x}_r(t) = f_r(\hat{x}_r(t), u(t), \hat{y}_r(t), \hat{\Sigma}_r(t)),
\]

such that

\[
\lim_{t \to \infty} (x(t) - \hat{x}_r(t)) = 0.
\]

In Problem 2, the messages received from the neighbor nodes by node \( i \) are \( \{\tilde{y}_{ij}(t), \tilde{\Sigma}_{ij}(t)\} \) \( (j \in J_i) \). Then, the all nodes estimate the state as follows:

**Problem 2.** Assume that \( (A, \tilde{C}_r) \) for all \( i \in V \) are the observable pairs. Given \( \Lambda_i \) for all \( i \in V \). Then, find the communication law

\[
\tilde{y}_{ij}(t) = \tilde{y}_{ij}(y_j, \{\tilde{y}_{ik}(t - D_{ik}); k \in J_i\}), \\
\tilde{\Sigma}_{ij}(t) = \tilde{h}_{ij}(\tilde{\Sigma}_{ik}(t - D_{ik}); j \in J_i), \\
\{u(\tau); t - D_{i,\max} \leq \tau \leq t\},
\]

and the distributed observers

\[
\tilde{y}_i(t) = \tilde{y}_i(y_j, \{\tilde{y}_{ik}(t - D_{ik}); k \in J_i\}), \\
\tilde{\Sigma}_i(t) = \tilde{h}_i(\tilde{\Sigma}_{ik}(t - D_{ik}); j \in J_i), \\
\{u(\tau); t - D_{i,\max} \leq \tau \leq t\},
\]

such that

\[
\lim_{t \to \infty} (x(t) - \tilde{x}_i(t)) = 0.
\]

We will obtain a result for Problem 1 in Sect. 3.2, and then extend it to Problem 2 in Sect. 3.3.

### 3. Proposed Method

#### 3.1 Preliminary

In this subsection, the past work which will be utilized for a delay compensation in the proposed method is introduced. The data received at each node includes multiple delays because the communication delays in the sensor network depend on the selection of communication paths. An observer with multi-sensor delays was proposed in [16]. Watanabe et al. [16] defines the outputs \( C_i x \) for each corresponding delay \( D_i \). Thus the output from all the measurements is expressed by

\[
y_{\text{cen}}(t) = \sum_{i=1}^{N} C_i x(t - D_i). \tag{6}
\]

We can compensate the delays in (6) by predicting the system behavior.

**Lemma 1** (Delay-Compensation Based on Prediction). Consider system (1) with output (6). For this system, the following equation holds:

\[
y_{\text{cen}}(t) = \sum_{i=1}^{N} C_i e^{-AD_i} \int_{t-D_i}^{t} e^{A(t-\tau)} B u(\tau) d\tau = \hat{C} x(t), \tag{7}
\]

where \( \hat{C} = \sum_{i=1}^{N} C_i e^{-AD_i}. \)
Proof. The solution for system (1) is expressed by
\[
x(t) = e^{AD}x(t - D_i) + \int_{t - D_i}^{t} e^{A(t - \tau)}Bu(\tau)d\tau.
\] (8)

Lemma 1 can be proven by solving (8) with respect to \(x(t - D_i)\) and inserting it into (6). \(\square\)

From Lemma 1, the state estimation for the system (1) with (6) becomes a finite pole assignment problem as follows.

**Lemma 2** (Watanabe’s Delay-Compensated Observer [16]). Consider system (1), the output (6), and the observer
\[
\hat{x}(t) = A\hat{x}(t) + Bu(t)
\]
\[
+ L\left(y_{cen}(t) + \sum_{i=1}^{N} C_i e^{-AD_i} \int_{t - D_i}^{t} e^{A(t - \tau)}Bu(\tau)d\tau - \sum_{i=1}^{N} C_i e^{-AD_i} \hat{x}(t)\right),
\] (9)

and suppose that \((A, \hat{\mathcal{C}})\) is an observable pair. Then, the estimation error \(\hat{x} = x - \hat{x}\) converges to zero, if and only if \(A - L\hat{\mathcal{C}}\) is Hurwitz.

**Proof.** From Lemma 1, the dynamics of \(\hat{x}\) can be expressed by
\[
\dot{\hat{x}} = (A - L\hat{\mathcal{C}})\hat{x}.
\] (10)

Therefore, \(\hat{x}\) tends to zero as \(t \to \infty\) if and only if \(A - L\hat{\mathcal{C}}\) is Hurwitz. \(\square\)

We notice that Watanabe’s delay-compensated observer does not need to handle \(y\), as defined by (2). The single output vector \(y_{cen}\) defined by (6), which includes all the delayed sensor signals, and the input signal \(u(\tau)\) \((t - \max(D_i)) \leq \tau \leq t\) are only required for the external signals of the observer (9). This property is effective for reducing the network traffic, because \(y\) includes redundant information. In addition, it is not assumed that the number of delay values \(N\) is smaller than dimension of the output \(m\). Thus, outputs that have the same elements but include different delays can be aggregated into one value. Based on these results, this paper solves Problems 1 and 2 in the following subsections.

### 3.2 Tree-Data-Aggregation-Based Observer

In this subsection, we consider Problem 1. Let \(r\) be a root node of \(T\). To collect information on the sensor network, each node executes the following communication. Let \(\hat{y}_i(t)\) be an aggregated output value at node \(i\). Each node aggregates its own measurements and the data received from child nodes, and sends these data to the parent node. The data sent from node \(i\) are expressed by
\[
\hat{y}_i(t) = y_i(t) + \sum_{j \in H_r^i} \hat{y}_j(t - D_{ij}).
\] (11)

Leaf node \(l\), which has no child node, i.e., \(H_r = \emptyset\), does not receive any data from the other nodes. Therefore, \(\hat{y}_i(t) = y_i(t)\) for each leaf node \(l\).

The aggregated measurements in each node include communication delays, which depend on the communication paths. To compensate these delays, each node calculates delay-compensation terms using the memory of the input, and sends the correction terms to the parent node. Let \(\hat{\mathcal{Z}}_i(t)\) be a variable that includes delay-compensation terms at node \(i\) and \(\mathcal{C}_i\) be a coefficient matrix that is recursively defined by
\[
\begin{align*}
\mathcal{C}_i &= C_i + \sum_{j \in H_r} \mathcal{C}_j e^{-AD_j} \quad (H_r \neq \emptyset), \\
\hat{\mathcal{C}}_i &= C_i \quad (H_r = \emptyset).
\end{align*}
\] (12)

Node \(i\) receives \(\hat{\mathcal{Z}}_j(t - D_{ij})\) from node \(j \in H_r\) and calculates \(\hat{\mathcal{Z}}_i(t)\) to compensate \(D_{ij}\) \((j \in H_r)\) as
\[
\hat{\mathcal{Z}}_i(t) = \sum_{j \in H_r} \left(\hat{\mathcal{Z}}_j(t - D_{ij}) + \mathcal{C}_j e^{-AD_j} \int_{t - D_{ij}}^{t} e^{A(t - \tau)}Bu(\tau)d\tau\right).
\] (13)

Each leaf node \(l\) does not need to calculate \(\hat{\mathcal{Z}}_i(t)\) because there are no data sent from the other nodes, which means that \(\hat{\mathcal{Z}}_i(t) = 0\) \((H_r = \emptyset)\).

Then, the following lemma holds for the communication law (11).

**Lemma 3** (Data Aggregation on Tree Networks). The aggregated value at the root node can be expressed by
\[
\hat{y}_r(t) = \sum_{i=1}^{N} C_i x(t - D_i),
\] (14)

which includes all the measurements on the network with delays.

**Proof.** Let \(\mathcal{H}_h^r\) be a set of nodes that can be reached from \(r\) via a simple path with length \(h\). It is expressed by
\[
\mathcal{H}_h^r = \begin{cases} \{r\} & (h = 0) \\ H_r & (h = 1) \\ \{i \in H_r; j \in \mathcal{H}_{h-1}^r\} & (h > 1). \end{cases}
\]

Moreover, we define
\[
\mathcal{H}_h^r = \bigcup_{i=0}^{h} \mathcal{H}_i^r.
\]

Using (11) twice, a relation
\[
\hat{y}_r(t) = \sum_{i \in \mathcal{H}_h^r} y_i(t - D_i) + \sum_{j \in \mathcal{H}_h^r} \hat{y}_j(t - D_{ij})
\] (15)

can be obtained. Because \(\hat{y}_r(t) = y_r(t)\) when node \(i\) satisfies \(H_r = \emptyset\), \(\hat{y}_r(t)\) is recursively given by (14). \(\square\)
By applying (12) twice,

**Proof.** By applying (12) twice,

\[
\check{C}_r = \sum_{i \in H_r} C_i e^{-AD_i} + \sum_{j \in H_r} \tilde{C}_j e^{-AD_j}
\]

(16)

is obtained. Note that \( \check{C}_i = C_i \) if \( H_r = \emptyset \). Therefore, matrix \( \check{C}_r \) is recursively given by

\[
\check{C}_r = \sum_{i=1}^{N} C_i e^{-AD_i}.
\]

The aggregated value \( \hat{\mathcal{Z}}_r(t) \) is given by (13). In addition, the data received from the child nodes are expressed as

\[
\hat{\mathcal{Z}}_j(t) = \sum_{k \in H_j} \left( \hat{\mathcal{Z}}_j(t - D_{jk}) + \tilde{C}_k e^{-AD_k} \int_{t-D_{jk}}^{t} e^{A(t-\tau)} Bu(\tau)d\tau \right).
\]

By substituting (18) in (13), we get

\[
\hat{\mathcal{Z}}_r(t) = \sum_{i \in H_r} C_i e^{-AD_i} \int_{t-D_i}^{t} e^{A(t-\tau)} Bu(\tau)d\tau
\]

\[
+ \sum_{j \in H_r} \left( \hat{\mathcal{Z}}_j(t - D_{ij}) + \tilde{C}_j e^{-AD_j} \int_{t-D_{ij}}^{t} e^{A(t-\tau)} Bu(\tau)d\tau \right).
\]

(19)

If node \( i \) is the leaf node, \( \hat{\mathcal{Z}}_i(t) = 0 \). Thus, \( \hat{\mathcal{Z}}_r(t) \) is recursively given by

\[
\hat{\mathcal{Z}}_r(t) = \sum_{i=1}^{N} C_i e^{-AD_i} \int_{t-D_i}^{t} e^{A(t-\tau)} Bu(\tau)d\tau.
\]

(20)

From (17), (20), and Lemmas 1 and 3, we can prove Lemma 4.

**Theorem 1.** Assume that \((A, \check{C}_r)\) is an observable pair. Then, the observer of the root node

\[
\dot{\hat{x}}_r(t) = A\hat{x}_r(t) + Bu(t) + L_r (\tilde{y}_r(t) - \hat{C}_r \hat{x}_r(t))
\]

(21)

can estimate the state, i.e., \( \hat{x}_r(t) \rightarrow x(t) \) as \( t \rightarrow \infty \), if and only if \( A - \check{C}_r L_r \) is Hurwitz.

**Proof.** Let \( \hat{x}_r(t) = x(t) - \hat{x}_r(t) \). From Lemma 4, the estimation error \( \hat{x}_r(t) \) satisfies the following equation:

\[
\dot{\hat{x}}_r(t) = (A - \check{C}_r L_r) \hat{x}_r(t).
\]

(22)

Thus, Theorem 1 is proven.

**3.3 Delay-Compensated Observer for Sensor Network**

In the previous subsection, the observed information are aggregated in the communication paths, and finally the root node can obtain an aggregated value for all the nodes’ information. However, with the exception of the root, all of the nodes only have part of the information observed by all the sensors. Each node needs the observed information of all the other nodes to estimate the state at the node. Because every node of a tree can be a root, each node can collect the observed values of all the other nodes in the same...
way as the method discussed in Sect. 3.2. However, waste-
ful communications will occur if we individually design the
communication laws to allow the different roots to collect
data. Therefore, in this subsection, we propose an efficient
intercommunication-based data aggregation method to es-
imate the state at all the nodes.

Let \( \bar{y}_{ij} \) and \( \bar{\Xi}_{ij} \) denote the data sent from node \( i \) to \( j \),
which will be defined later. The data sent from node \( i \) to \( j \)
are the aggregated information from the neighbor nodes of
node \( i \), except for \( j \), i.e., \( J_i \setminus \{ j \} \). Therefore, \( \bar{y}_{ij} \) and \( \bar{\Xi}_{ij} \) can be defined by

\[
\bar{y}_{ij}(t) = y_i(t) + \sum_{k \in J_i \setminus \{ j \}} \bar{y}_{ki}(t - D_{ik}),
\]

\[
\bar{\Xi}_{ij}(t) = \sum_{k \in J_i \setminus \{ j \}} \left( \bar{\Xi}_{ki}(t - D_{ik}) + \tilde{C}_{ki} e^{-AD_{ik}} \int_{t-D_{ik}}^{t} e^{A(t-\tau)} B u(\tau) d\tau \right),
\]

where \( \tilde{C}_{ij} \) is the matrix expressed by

\[
\tilde{C}_{ij} = C_i + \sum_{k \in J_i} \tilde{C}_{ik} e^{-AD_{ik}}.
\]

Note that matrix \( \tilde{C}_{ij} \) can be obtained by an offline calculation.
The aggregated values of each node, \( \bar{y}_{ij}(t) \) and \( \bar{\Xi}_{ij}(t) \), are given by

\[
\bar{y}_{ij}(t) = y_i(t) + \sum_{j \in J_i} \bar{y}_{ji}(t - D_{ij}),
\]

\[
\bar{\Xi}_{ij}(t) = \sum_{j \in J_i} \left( \bar{\Xi}_{ji}(t) + \tilde{C}_{ji} e^{-AD_{ji}} \int_{t-D_{ji}}^{t} e^{A(t-\tau)} B u(\tau) d\tau \right),
\]

and the output matrix after the delay-compensation is recu-
rsively defined by

\[
\hat{C}_i = C_i + \sum_{j \in J_i} \tilde{C}_{ij} e^{-AD_{ij}}.
\]

Using the intercommunication law in (23) and (24), the fol-
lowing theorem holds.

**Theorem 2.** The aggregated value of each node \( \bar{y}_i(t) \) in-
cludes the outputs of all the sensors. The delays included in
\( \bar{y}_i(t) \) can be compensated by \( \bar{\Xi}_i(t) \). Let \( \hat{x}_i \) be an estimate
of the state calculated at node \( i \). The observer in node \( i \) is defined by

\[
\hat{x}_i(t) = A \hat{x}_i(t) + B u(t) + L_i \left( \bar{y}_i(t) + \bar{\Xi}_i(t) - \hat{C}_i \hat{x}_i(t) \right),
\]

where \( (A, \hat{C}_i) \) is an observable pair. Then, the error dyna-
mics of (29) for node \( i \) are asymptotically stable if \( A - L_i \hat{C}_i \) is
Hurwitz.

**Proof.** The intercommunication law of (23) and (24) im-
plies that \( \bar{y}_{ij}(t) \) and \( \bar{\Xi}_{ij}(t) \) are equal to the aggregated data
expressed by (11) and (13), respectively, when node \( j \) is the
parent node of node \( i \). Similarly, \( \hat{C}_{ij} \) in (25) coincides with
\( \tilde{C}_i \) of (12), when the parent node is \( j \). Thus, the aggregated
values of node \( i \), which are defined in (26) and (27), become
the tree-based aggregated values at the root node. Matrix
(28) also becomes the aggregated matrix whose root node is
its own node. Therefore, the error dynamics of the observer
of each node (29) are asymptotically stable if \( A - \hat{C}_i L_i \) is
Hurwitz.

\[ \square \]

**Remark 3** (Observability of Sensor Network). In Theorem
2, it is assumed that all the pairs \((A, \hat{C}_i)\) are observable.
Thus, the observability of the sensor network depends on
the network topology because \( \hat{C}_i \) includes the delay values.
In general, the condition that pair \( A \) and non-delay output matrix \( C_{\text{avg}} \) are observable does not guarantee that \((A, \hat{C}_i)\) is observable. However, we can expect that the sensor net-
work becomes observable if \((A, C_{\text{avg}})\) is observable and the
communication delays are sufficiently small.

### 3.4 Adaptation Algorithm for Modification of Network Topology

In this subsection, we show an algorithm for the recalcula-
tion of the parameters in each node when the topology of
the network is modified. The parameters that depend on the
network topology in each node are \( \hat{C}_i, \hat{C}_{ij}, \) and \( L_i \). The pro-
posed algorithm calculates these parameters through local
calculations and mutually communications between nodes.

We define logical variables \( \delta_i(t) \) as \( \delta_i(t) \in \{T,F\} \),
which indicates whether \( J_i \) is modified. Node \( i \) sets \( \delta_i(t) \) to ‘T’
(true) if \( J_i \) has been modified at \( t \), and otherwise it \( \delta_i(t) = F \)
(false). The signals \( \delta_i(t) \) represent the propagation of the
modification from node \( i \). If node \( i \) needs to tell the present
of the modification to node \( j \), \( \delta_{ij}(t) \) becomes ‘T,’ which
means

\[ \delta_{ij}(t) = \delta_i(t) \lor \left( \bigvee_{k \in J_i \setminus \{ j \}} \delta_{ki}(t - D_{ik}) \right). \]

According to \( \delta_i(t) \) and \( \delta_{ij}(t) \), each node recalculates or
updates each parameter. If \( \delta_{ij}(t) \) is ‘T,’ node \( i \) executes an
event to update \( \hat{C}_{ij} \) based on (25). Let \( \delta_i(t) \) be

\[
\delta_i(t) = \delta_i(t) \lor \left( \bigvee_{j \in J_i} \delta_{ij}(t - D_{ij}) \right).
\]

Node \( i \) needs to recalculate \( L_i \) and \( \hat{C}_i \) when \( \delta_i(t) \) is ‘T.’ The
triggered node updates \( \hat{C}_i \) as (28), and chooses \( L_i \) such that
\( A - L_i \hat{C}_i \) becomes stable. Algorithm 1 have summarized the
above procedure.

To execute Algorithm 1, each node needs to prepare \( C_i \)
and \( e^{-AD_{ij}} \) for all \( j \) which are candidates for the neighbor-
hood nodes. Each node can utilize unsteady Kalman filter
algorithms to recalculate $L_i$. The unsteady Kalman filter algorithms need to calculate Riccati differential equation, but do not need the real-time calculation of eigenvalues or an inverse matrix. Therefore, the assumption that each node has the ability to execute Algorithm 1 is reasonable.

4. Numerical Simulation

Let us consider the quadruple tank system \[21\] as follows:

$$A = \begin{pmatrix}
-0.15 & 0 & 0.5 & 0 \\
0 & -0.25 & 0 & 0.5 \\
0 & 0 & -0.15 & 0 \\
0 & 0 & 0 & -0.25
\end{pmatrix},$$

$$B = \begin{pmatrix}
0.4 & 0 & 0 & 0.4 \\
0 & 0.4 & 0 & 0
\end{pmatrix}^T,$$

$$u = \left(\cos(\pi t) + 1, \sin(\pi t) + 1\right)^T.$$

There are 50 sensor nodes with the communication paths which are illustrated in Fig. 4. Each sensor in Fig. 4 measures $x_1$ or $x_2$. The topology of the sensor network is generated by a BA model. Let $v_i := (x_{gi}, y_{gi})$ be a coordinate of node $i$ in Fig. 4. We set the observation matrix of each node as

Algorithm 1 Recalculation of $\hat{C}_i$, $\bar{C}_{ij}$ and $L_i$ in Each Time Sequence

if $J_i$ is modified at $t$ then
  $\delta_i(t) \leftarrow$ true
else
  $\delta_i(t) \leftarrow$ false
end if

for all $j$ such that $j \in J_i$ do
  if $\bar{\delta}_{ji}(t-D_{ij}) = \text{true}$ then
    $\bar{C}_{ji} \leftarrow \bar{C}_{ji}^{\text{new}}$
  end if
end for

for all $j$ such that $j \in J_i$ do
  $\bar{\delta}_{ij}(t) \leftarrow \bar{\delta}_{ij}(t) \lor \left(\bigvee_{k \in J_i} \bar{\delta}_{ki}(t-D_{jk})\right)$
  if $\bar{\delta}_{ij}(t) = \text{true}$ then
    $\bar{C}_{ij}^{\text{new}} \leftarrow C_i + \sum_{k \in J_i} \bar{C}_{ik} e^{-AD_{jk}}$
  end if
end for

$\hat{\delta}_i(t) \leftarrow \hat{\delta}_i(t) \lor \left(\bigvee_{k \in J_i} \bar{\delta}_{ki}(t-D_{jk})\right)$

if $\hat{\delta}_i(t) = \text{true}$ then
  $\hat{C}_i \leftarrow C_i + \sum_{k \in J_i} \bar{C}_{ik} e^{-AD_{jk}}$
end if

Calculate $L_i$ such that $A - L_i \hat{C}_i$ becomes stable.

$$C_i = \begin{cases}
0.02 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 \\
0.02 & 0 & 0 & 0
\end{cases},$$

if $x_{gi} > 0$ otherwise.
The communication delays on \((i, j) \in E\) are given by \(D_{ij} = |e_i - e_j| \times 0.1\).

In a first simulation, the three topologies in Figs. 5, 6 and 7 are switched every 4 seconds. These topologies are the subgraphs of Fig. 4 which have tree structures. Figure 8 shows the time responses of the estimation errors on \(x_1\), \(x_2\), \(x_3\), and \(x_4\). After the estimation errors converge to zero, the parameter modifications do not affect the estimates. Therefore, we can confirm that the observer proposed in this paper can estimate the state over the switched tree. In a second simulation, Fig. 5 with communication jitters are used. Let the delays with the jitters be

\[
\hat{D}_{ij} = (\xi + 1)D_{ij},
\]

where \(\xi \in [0, 0.15]\) is the uniform random number. Figure 9 shows the time responses of the estimation errors on \(x_1\), \(x_2\), \(x_3\), and \(x_4\) with the communication jitters. The proposed method can estimate the states with a sufficient accuracy over the network with the communication jitters. Therefore, these numerical simulation results verify the effectiveness of the proposed observer and data aggregation method.

5. Conclusion

This paper proposes a data-aggregation-based delay compensated observer for a wireless sensor network. The proposed method aggregates the values measured by all the sensors to each node. The communication delay between the neighbor nodes is instantly compensated by each node. Therefore, all of the nodes of a sensor network can estimate the state of the system. The dimensions of the signals on the communication paths is \(2n\), which is independent of node number \(N\). This implies that the proposed communication laws are scalable with respect to the network size. A numerical simulation verifies the effectiveness of the proposed method.

In the proposed method, it is assumed that all the nodes can obtain inputs in real time via a fast network. To remove this assumption, we will consider an observer-based distributed controller in a future study. Moreover, the estimated state of the proposed method and output of the sensor networks have redundancy. We believe that this redundancy could enable us to realize a fault tolerant design for the distributed observers or controllers in wireless networks. This will also be the focus of our future study.

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